



Imperial College London

Accommodating fairness in a shared-energy allocation problem with uncertainties ICSP 2023 - UC Davis

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WHAT IS A PROSUMER?



AN EXAMPLE OF ENERGY MARKET



AN EXAMPLE OF ENERGY MARKET





A NEED FOR AGGREGATION



A NEED FOR AGGREGATION



A NEED FOR AGGREGATION



➡ How to fairly allocate benefits?

. time

1 Defining fairness

2 Modeling a fair prosumer aggregation

B Dynamic extension

4 Stochastic extension

5 Conclusion

Definition (Oxford Dictionary)

Fairness is the quality of treating people equally or in a way that is reasonable.

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Fairness is the quality of treating people equally or in a way that is reasonable.

- A take on fairness is necessarily subjective, as the definition is ambiguous.
- ► Moreover, are we looking for fair outcomes or for fair processes?
- ► Mathematic modeling requires a clear definition, leading to a particular model.
- Then we must make a choice which has consequences on the solution.

- Egalitarism: everyone gets the same share.
- Rawl's theory (minimax): we should favor the least well-off.
- The Need Principle: first, we satisfy everyone's basic needs, then we can focus on efficiency.
- **Proportional fairness**: derived from Nash bargaining solution in Game theory, the allocation must satisfy some properties.

⇒ scale invariance, symmetry, Pareto optimality, independence of irrelevant alternatives



Ice Cream individual 6\$

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Individualist



Ice Creams pack 31\$





Ice Cream individual 6\$











- The literature tackles fairness considerations in various fields.
 - game theory, communications networks, facility locations, portfolio optimization, machine learning . . .

FAIRNESS BY DESIGN

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- A natural approach is to solve a problem efficiently and then allocate costs fairly.
 - computing Shapley's values
 - having allocation policies

FAIRNESS BY DESIGN

- The literature tackles fairness considerations in various fields.
 - game theory, communications networks, facility locations, portfolio optimization, machine learning . . .
- A natural approach is to solve a problem efficiently and then allocate costs fairly.
 - computing Shapley's values
 - having allocation policies
- However, we focus on fairness by design: fairness is accommodated in the model.
 - In A Guide to Formulating Equity and Fairness in an Optimization Model, Violet (Xinying) Chen and J. N. Hooker review how to model fairness in an optimization model.

Defining fairness

2 Modeling a fair prosumer aggregation

Modeling a prosumer Modeling a prosumer aggregation

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TOY EXAMPLE: CONTEXT



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MODEL FOR ONE PROSUMER

Prosumer model $(A_i) := \min_{y^{A_i}, b} f_i(y^{A_i}, b)$

- y^{A_i} are the control variables (energy purchases) for A_i
- *b* represents the decision to buy in advance or not
- Minimize energy costs

$$f_i(y^{A_i}, b) = \sum_{t \in [T]} p_t^{\mathsf{DA}} q_{t,\mathsf{DA}}^{A_i} + p_t^{\mathsf{B}} q_{t,\mathsf{B}}^{A_i}$$

MODEL FOR ONE PROSUMER

Prosumer model

$$(A_i) := \underset{\substack{y^{A_i}, b}{\text{s.t}}}{\min} f_i(y^{A_i}, b)$$
s.t $y^{A_i} \in \mathcal{Y}^{A_i},$

$$\mathcal{Y}^{A_i} := \begin{cases} \frac{Q_i \leq q_{t,\mathsf{DA}}^{A_i} + q_{t,\mathsf{B}}^{A_i} \leq \overline{Q_i}, \ t \in [T] \text{ bounded load} \\ \sum_{t=1}^{T} (q_{t,\mathsf{DA}}^{A_i} + q_{t,\mathsf{B}}^{A_i}) \geq L^i \text{ overall load requirements} \end{cases}$$

MODEL FOR ONE PROSUMER

Prosumer model
$$(A_i) :=$$
 $\min_{y^{A_i}, b}$ $f_i(y^{A_i}, b)$ s.t $y^{A_i} \in \mathcal{Y}^{A_i},$ $b \in \mathcal{B}(y^{A_i}).$

$$\mathcal{B}(y^{A_i}) = \left\{ \underline{q_{t,\mathsf{DA}}} b_t \le q_{t,\mathsf{DA}}^{A_i} \le M b_t, \ b_t \in \{0,1\}, \ t \in [T] \right\}$$
minimum day-ahead purchases

MODEL FOR A PROSUMER AGGREGATION

Aggregator modeli	ng	
$(A) := \min_{y,b}$	f(y, b)	
s.t	$y^{A_i} \in \mathcal{Y}^{A_i}$ $b \in \mathcal{B}(y)$	$\forall i$

• *f* is the objective function

• $y := (y^{A_1}, \dots, y^{A_N})$ are the control variables for all prosumer

$$\mathcal{B}(y) = \left\{ \underline{q_{t,\mathsf{DA}}} b_t \leq \sum_{i=1}^{N} q_{t,\mathsf{DA}}^{A_i} \leq M b_t, \ b_t \in \{0,1\}, \ t \in [T] \right\}$$

MODEL FOR A PROSUMER AGGREGATION

Aggregator modeling
$$(A) := Min_{y,b} f(y,b)$$
s.t $y^{A_i} \in \mathcal{Y}^{A_i}$ $b \in \mathcal{B}(y)$

We have two challenges with the aggregation :

1. Acceptability

• add constraints $c_{A_i} \leq \overline{c}_{A_i}$

where
$$c_{A_i} := \sum_{t=1}^{T} c_{A_i,t}$$
 is the cost of A_i in (A),
and $\overline{c}_{A_i} := \sum_{t=1}^{T} \overline{c}_{A_i,t}$ is the optimal value of (A_i) .

MODEL FOR A PROSUMER AGGREGATION

Aggregator modeling
$$(A) := Min_{y,b} f(y,b)$$
s.t $y^{A_i} \in \mathcal{Y}^{A_i}$ $b \in \mathcal{B}(y)$

We have two challenges with the aggregation :

1. Acceptability

2. Fairness

• Utilitarian:
$$f(y) = \sum_{i=1}^{N} c_{A_i}$$
.

• Minimax proportional:
$$f(y) := \underset{i \in [\![1,N]\!]}{\mathsf{Max}} \frac{\overline{c}_{A_i} - c_{A_i}}{\overline{c}_{A_i}}.$$

• Proportional:
$$f(y) := \eta \sum_{i=1}^{N} \log(c_{A_i}).$$

TOY EXAMPLE: DATA





TOY EXAMPLE: INDIVIDUALIST





TOY EXAMPLE: UTILITARIAN





	CA_1	CA2	$c_{A_1} + c_{A_2}$
Individualist	250	50	300
Utilitarian	130	52	182
Acceptable Utilitarian			
Proportional			
Minimax proportional			

TOY EXAMPLE: ACCEPTABLE UTILITARIAN

Day-ahead (\$/MWh) Balancing (\$/MWh)





	CA1	CA ₂	$c_{A_1} + c_{A_2}$
Individualist	250	50	300
Utilitarian	130	52	182
Acceptable Utilitarian	155	46	201
Proportional			
Minimax proportional			

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TOY EXAMPLE: PROPORTIONAL





	CA1	CA_2	$c_{A_1} + c_{A_2}$
Individualist	250	50	300
Utilitarian	130	52	182
Acceptable Utilitarian	155	46	201
Proportional	220	20	240
Minimax proportional			

TOY EXAMPLE: MINIMAX PROPORTIONAL

Day-ahead (\$/MWh) Balancing (\$/MWh)





	CA1	CA2	$c_{A_1} + c_{A_2}$
Individualist	250	50	300
Utilitarian	130	52	182
Acceptable Utilitarian	155	46	201
Proportional	220	20	240
Minimax proportional	170	32	202

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- The utilitarian is the most efficient approach, but it can lead to unacceptable solutions for some prosumers.
- The proportional approach coming from a theory of bargaining favors smaller prosumers.
- The minimax approach derived from Rawl's theory on fairness: the solution finds a balance between both prosumers.

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With long term problems in mind, we extend the previous model to the dynamic case.

We adapt acceptability constraint so that prosumers have no incentive to leave the aggregation between stages.

EXTENTION: THE DYNAMIC CASE

With long term problems in mind, we extend the previous model to the dynamic case.

We adapt acceptability constraint so that prosumers have no incentive to leave the aggregation between stages.

• Average Acceptability:

$$\sum_{t=1}^{T} c_{A_{i},t} \leq \sum_{t=1}^{T} \overline{c}_{A_{i},t}, \qquad \forall i$$

• Progressive Acceptability:

$$\sum_{ au=1}^t c_{\mathcal{A}_i, au} \leq \sum_{ au=1}^t \overline{c}_{\mathcal{A}_i, au}, \qquad orall i, orall t$$

• Stage-wise Acceptability:

$$c_{A_i,t} \leq \overline{c}_{A_i,t}, \quad \forall i, \forall t$$

TOY EXAMPLE: DYNAMIC CASE



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TOY EXAMPLE: AVERAGE ACCEPTABILITY





	CA_1	CA2	$c_{A_1} + c_{A_2}$
Individualist	250	50	300
Utilitarian	130	52	182
Average			0.01
Acceptability	155	46	201
Progressive			
Acceptability			
Stage-wise			
Acceptability			

TOY EXAMPLE: PROGRESSIVE ACCEPTABILITY





CA1	CA_2	$c_{A_1} + c_{A_2}$
250	50	300
130	52	182
155	10	001
155	46	201
170	20	000
170	32	202
	cA1 250 130 155 170	$ \begin{array}{c cccc} c_{A_1} & c_{A_2} \\ \hline 250 & 50 \\ 130 & 52 \\ \hline 155 & 46 \\ \hline 170 & 32 \\ \hline \end{array} $

TOY EXAMPLE: STAGE-WISE ACCEPTABILITY





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 We now consider that the balancing prices {p_t^B}_{t∈[T]} are random variables.

STOCHASTIC EXTENSION

- We now consider that the balancing prices {p_t^B}_{t∈[T]} are random variables.
- In a 2 stage setting, we have:
 - 1. first-stage variables: day-ahead purchases $\{b_t, q_{t,DA}^i\}_{t \in [T], i \in [N]}$
 - 2. recourse variables: balancing purchases $\{q_{t,B}^i\}_{t \in [T], i \in [N]}$



1. How to handle the uncertainties?

• Optimize expected costs:

 $\underset{q_{DA},b,q_{B}}{\mathsf{Min}} \quad \mathbb{E}_{\omega} \left[f(q_{DA}, b, q_{B}, \omega) \right]$

• Optimize the worst case:

 $\underset{q_{DA},b,q_{B}}{\text{Min}} \quad \underset{\omega \in \Omega}{\text{Max}} \left\{ f(q_{DA},b,q_{B},\omega) \right\}$

• Optimize a risk measure:

 $\underset{q_{\mathsf{,DA}}, b, q_{\mathsf{,B}}}{\mathsf{Min}} \quad \rho_{\omega} \left\{ f(q_{\mathsf{,DA}}, b, q_{\mathsf{,B}}, \omega) \right\}$

- 1. How to handle the uncertainties?
- 2. How to accommodate fairness into the model?

• Utilitarian:

$$f(y) = \sum_{i=1}^{N} c_{A_i}$$

• Minimax proportional:

$$f(y) := \max_{i \in \llbracket 1, N
rbracket} rac{\overline{c}_{\mathcal{A}_i} - c_{\mathcal{A}_i}}{\overline{c}_{\mathcal{A}_i}}.$$

- 1. How to handle the uncertainties?
- 2. How to accommodate fairness into the model?
- 3. How to adapt the acceptability constraints?

• Expected Acceptability:

$$\boldsymbol{c}_{\boldsymbol{\mathcal{A}}_i} \preceq_{\mathbb{E}} \overline{\boldsymbol{c}}_{\boldsymbol{\mathcal{A}}_i} \iff \mathbb{E}[\boldsymbol{c}_{\boldsymbol{\mathcal{A}}_i}] \leq \mathbb{E}[\overline{\boldsymbol{c}}_{\boldsymbol{\mathcal{A}}_i}] \qquad \forall i$$

ACCEPTABILITY CONSTRAINTS

• Expected Acceptability:

$$c_{A_i} \preceq_{\mathbb{E}} \overline{c}_{A_i} \iff \mathbb{E}[c_{A_i}] \leq \mathbb{E}[\overline{c}_{A_i}] \qquad \forall i$$

• Almost sure Acceptability:

$$c_{\mathcal{A}_i} \preceq_{as} \overline{c}_{\mathcal{A}_i} \iff c_{\mathcal{A}_i,\omega} \le \overline{c}_{\mathcal{A}_i,\omega}, \quad \forall \omega \qquad \forall i$$

ACCEPTABILITY CONSTRAINTS

• Expected Acceptability:

$$\boldsymbol{c}_{\boldsymbol{A}_i} \preceq_{\mathbb{E}} \overline{\boldsymbol{c}}_{\boldsymbol{A}_i} \iff \mathbb{E}[\boldsymbol{c}_{\boldsymbol{A}_i}] \leq \mathbb{E}[\overline{\boldsymbol{c}}_{\boldsymbol{A}_i}] \qquad \forall i$$

• Almost sure Acceptability:

$$c_{\mathcal{A}_i} \preceq_{as} \overline{c}_{\mathcal{A}_i} \iff c_{\mathcal{A}_i,\omega} \le \overline{c}_{\mathcal{A}_i,\omega}, \quad \forall \omega \qquad \forall i$$

• 1storder Acceptability:

$$c_{\mathcal{A}_i} \preceq_{(1)} \overline{c}_{\mathcal{A}_i} \iff \mathbb{P}(\overline{c}_{\mathcal{A}_i} \le \eta) \le \mathbb{P}(c_{\mathcal{A}_i} \le \eta), \ \forall \eta \in \mathbb{R} \qquad \forall i$$

• Expected Acceptability:

$$\boldsymbol{c}_{\boldsymbol{A}_i} \preceq_{\mathbb{E}} \overline{\boldsymbol{c}}_{\boldsymbol{A}_i} \iff \mathbb{E}[\boldsymbol{c}_{\boldsymbol{A}_i}] \leq \mathbb{E}[\overline{\boldsymbol{c}}_{\boldsymbol{A}_i}] \qquad \forall i$$

• Almost sure Acceptability:

$$c_{\mathcal{A}_i} \preceq_{as} \overline{c}_{\mathcal{A}_i} \iff c_{\mathcal{A}_i,\omega} \leq \overline{c}_{\mathcal{A}_i,\omega}, \quad \forall \omega \qquad \forall i$$

• 1storder Acceptability:

$$\boldsymbol{c}_{\boldsymbol{A}_{i}} \preceq_{(1)} \overline{\boldsymbol{c}}_{\boldsymbol{A}_{i}} \iff \mathbb{P}(\overline{\boldsymbol{c}}_{\boldsymbol{A}_{i}} \leq \eta) \leq \mathbb{P}(\boldsymbol{c}_{\boldsymbol{A}_{i}} \leq \eta), \ \forall \eta \in \mathbb{R} \qquad \forall i$$

• Increasing convex Acceptability:

$$\boldsymbol{c}_{\boldsymbol{A}_{i}} \leq_{ic} \overline{\boldsymbol{c}}_{\boldsymbol{A}_{i}} \iff \mathbb{E}\left[(\boldsymbol{c}_{\boldsymbol{A}_{i}}-\eta)^{+}\right] \leq \mathbb{E}\left[(\overline{\boldsymbol{c}}_{\boldsymbol{A}_{i}}-\eta)^{+}\right], \ \forall \eta \in \mathbb{R} \qquad \forall i$$

TOY EXAMPLE

With a similar example as the one presented in the deterministic case, we draw 100 scenarios for the balancing prices.

We test the small example with:

- different objective functions:
 - expected utilitarian costs
 - expected minimax costs

- robust utilitarian costs
- robust minimax costs
- the different acceptability constraints

TOY EXAMPLE: EXPECTED UTILITARIAN



TOY EXAMPLE: ALL RESULTS



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5 Conclusion

- Aggregation for prosumers saves costs, the question is how to fairly allocate them?
- We present two challenges in the problem:
 - 1. acceptability modeled through constraints;
 - 2. fair allocation modeled through the objective function.
- We extend fairness to dynamic and stochastic frameworks:
 - ▶ fairness can be derived to take into consideration cost distribution and risk aversion.